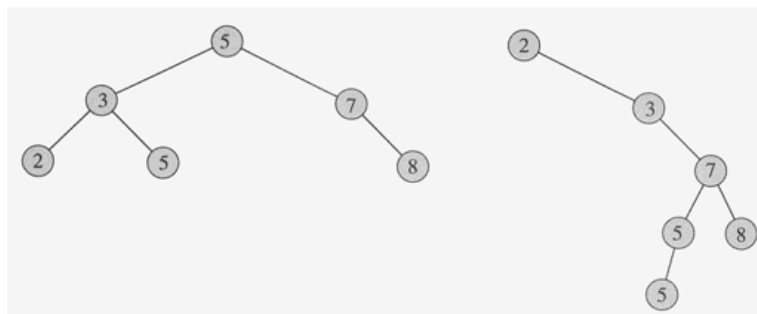


Binary Search Trees

Chapter 12

What is a Binary Search Tree?

- A binary tree where each node is an object
 - Each node has a key value, left child, and right child (might be empty)
- Each node satisfies the binary search tree property
 - Let x be a node in the BST. The left child's key must be $\leq x$'s key. The right child's key must be $\geq x$'s key



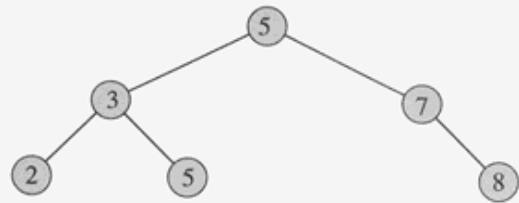
Traversing the BST

INORDER-TREE-WALK(x)

```

1  if  $x \neq \text{NIL}$ 
2      then INORDER-TREE-WALK( $\text{left}[x]$ )
3          print  $\text{key}[x]$ 
4      INORDER-TREE-WALK( $\text{right}[x]$ )

```



$O(n)$ time

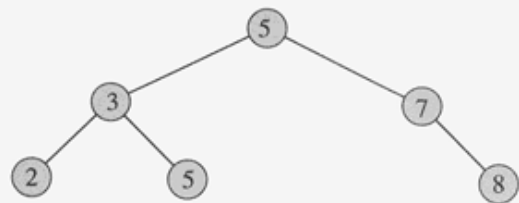
Searching a BST

TREE-SEARCH(x, k)

```

1  if  $x = \text{NIL}$  or  $k = \text{key}[x]$ 
2      then return  $x$ 
3  if  $k < \text{key}[x]$ 
4      then return TREE-SEARCH( $\text{left}[x], k$ )
5  else return TREE-SEARCH( $\text{right}[x], k$ )

```



Runs in $O(h)$ time but this could be $O(n)$ in the worst case!
 $O(\lg n)$ if the tree is balanced!

Finding min and max?

Successor

- Finding the node with the next largest (or equal) value

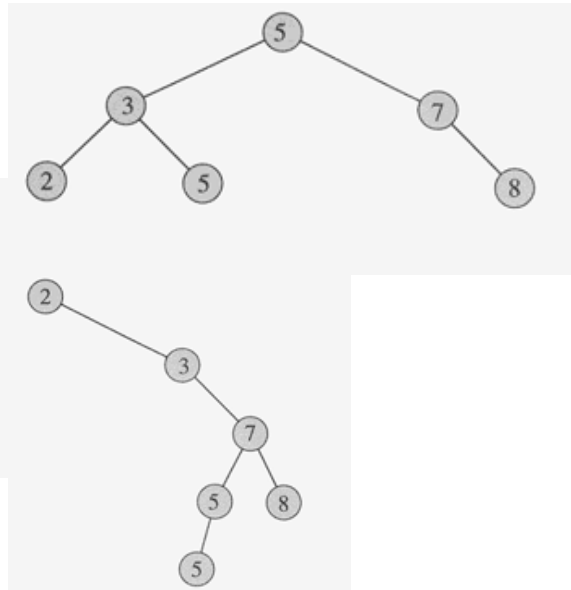
TREE-SUCCESSOR(x)

```

1  if  $right[x] \neq \text{NIL}$ 
2    then return TREE-MINIMUM( $right[x]$ )
3   $y \leftarrow p[x]$ 
4  while  $y \neq \text{NIL}$  and  $x = right[y]$ 
5    do  $x \leftarrow y$ 
6     $y \leftarrow p[y]$ 
7  return  $y$ 

```

$O(h)$ runtime



Insertion

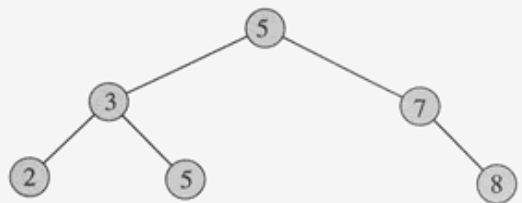
TREE-INSERT(T, z)

```

1   $y \leftarrow \text{NIL}$ 
2   $x \leftarrow \text{root}[T]$ 
3  while  $x \neq \text{NIL}$ 
4    do  $y \leftarrow x$ 
5    if  $key[z] < key[x]$ 
6      then  $x \leftarrow left[x]$ 
7    else  $x \leftarrow right[x]$ 
8   $p[z] \leftarrow y$ 
9  if  $y = \text{NIL}$ 
10   then  $\text{root}[T] \leftarrow z$ 
11   else if  $key[z] < key[y]$ 
12     then  $left[y] \leftarrow z$ 
13     else  $right[y] \leftarrow z$ 

```

▷ Tree T was empty

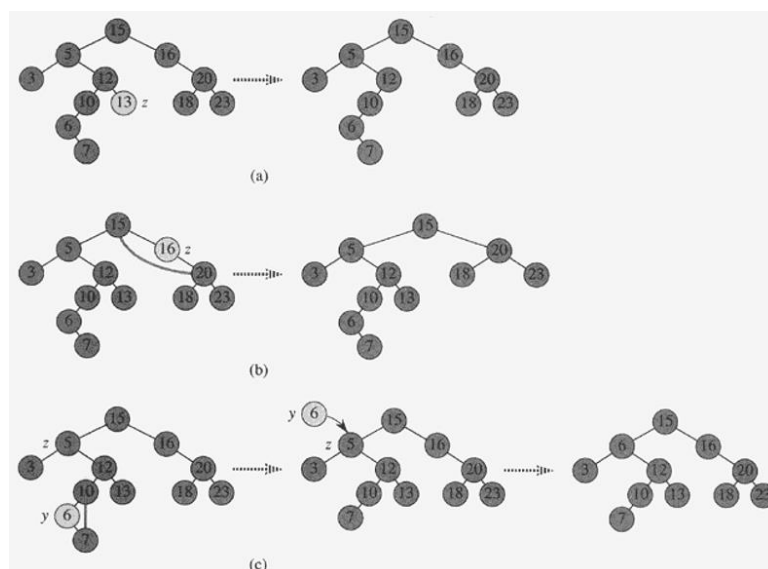


$O(h)$ runtime

Deletion

- Deleting a node z from a BST T
 1. If z has no children then simply remove it by modifying its parent to replace z with nil as its child
 2. If z has just one child then we elevate that child to take z 's position in the tree by modifying z 's parent to replace z by z 's child
 3. If z has two children then:
 - Find z 's successor y – which must be in z 's right subtree – and have y take z 's position in the tree
 - As a successor y in the right subtree, y has at most one child. Remove y using rule 2
 - The rest of z 's original right subtree becomes y 's right subtree and z 's left subtree becomes y 's left subtree

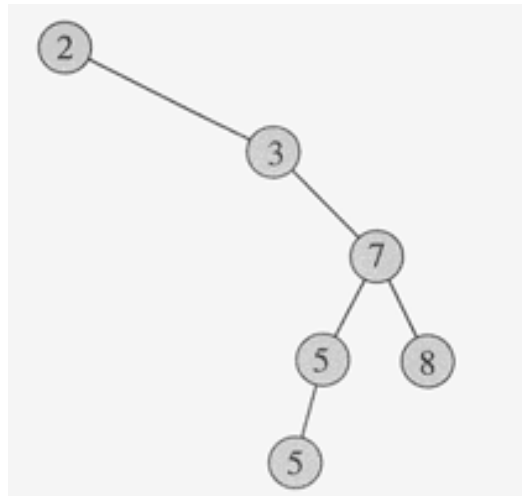
Delete Examples



Deletion

```

Tree-Delete(T,z)
if z.left == NIL
    Transplant(T, z, z.right)
elseif z.right == NIL
    Transplant(T, z, z.left)
else
    y = Tree-Minimum(z.right)
    if y.p != z
        Transplant(T,y,y.right)
        y.right = z.right
        y.right.p = y
    Transplant(T, z, y)
    y.left = z.left
    y.left.p = y
  
```



BST

- Worst case?
- Best case?
- Expectation for randomly built BST?